

## Knuth Equivalence + Green's Theorem

Def:  $\pi, \sigma \in S_n$  are said to be  $P$ -equivalent, written  $\pi \stackrel{P}{\cong} \sigma$ , if  $P(\pi) = P(\sigma)$ .

On  $S_3$  we have:

$\{123\}$        $\{213, 231\}$        $\{132, 312\}$        $\{321\}$

Def: Suppose  $x < y < z$ . Then if  $\pi, \sigma \in S_n$ , say  $\pi \stackrel{1}{\cong} \sigma$  if

1.  $\pi = x_1 \dots yxz \dots x_n$  and  $\sigma = x_1 \dots yzx \dots x_n$   
and  $\pi \stackrel{2}{\cong} \sigma$  if

2.  $\pi = x_1 \dots xzy \dots x_n$  and  $\sigma = x_1 \dots zxy \dots x_n$

Relation  $\stackrel{K}{\cong}$  is trans. closure of  $\stackrel{1}{\cong}$  and  $\stackrel{2}{\cong}$

Theorem (Knuth): If  $\pi, \sigma \in S_n$  then

$$\pi \stackrel{K}{\cong} \sigma \text{ iff } \pi \stackrel{P}{\cong} \sigma$$

Proof: ( $\Rightarrow$ ) Last time we saw enough to show that if  $\pi \stackrel{1}{\cong} \sigma$  then  $\pi \stackrel{P}{\cong} \sigma$

Now assume  $\pi \stackrel{2}{\cong} \sigma$ .

$$\pi = x_1 \dots yxz \dots x_n \text{ and } \sigma = x_1 \dots yzx \dots x_n$$

Let  $P$  be the tableau obtained by inserting the elements before  $y$ .

Need to show:  $r_z r_x r_y(P) = r_x r_z r_y(P)$ .

Idea: the insertion path of  $y$  creates a "barrier" s.t. paths for  $x$  and  $z$  lie to the left & right of this barrier, no matter in which order they are inserted. So since the paths don't intersect,  $r_x r_z + r_z r_x$  have same effect.

Induction on # of rows of  $P$ . If  $P$  has no rows, then  $P(zxy) = \begin{matrix} x & y \\ z \end{matrix} = P(xyz)$

Now let  $P$  have  $r > 0$  rows & consider  $\overline{P} = r_y(P)$ . Let  $k$  be column where  $y$  enters  $P$ , bumping element  $y'$ . Then:

$$\#1 \quad \overline{P}_{1,j} \leq y' \text{ for all } j \leq k$$

$$\#2 \quad \overline{P}_{1,l} > y' \text{ for all } l > k.$$

If we insert  $x$  next, since  $x < y$ ,  $x$  will enter in a column  $j$  w/  $j \leq k$ .

And if  $x'$  is bumped, then  $x' < y'$  by #1.

If we insert  $z$  into  $r_x r_y(P)$  then since  $z > y$ , by #1,  $z$  enters column  $l$  for  $l > k$ .

The bumped element  $z'$  satisfies  $z' > y'$  by #2.

Now consider  $r_x r_z r_y(P)$ . We have #1 and #2.

Again,  $z > y \Rightarrow$  by #1,  $z$  enters column  $l$  for  $l > k$ , bumping the same  $z'$  as above.

And  $x < y \Rightarrow x$  enters in a column  $j$  w/  $j \leq k$ , bumping the same  $x'$  as above. Note:  $x' < y' < z'$ .

So in both cases,  $z$  and  $x$  enter in columns strictly to right & weakly to left of column  $k$  — since columns are disjoint, insertions of  $x$  &  $z$  don't affect each other. So 1st rows of  $r_z r_x r_y(P)$  and  $r_x r_z r_y(P)$  are equal.

In both cases we displaced same elements  $x', y', z'$  w/  $x' < y' < z'$ , so by induction, the remainders of the two tableaux are equal. This proves  $(\Rightarrow)$ .

Before  $(\Leftarrow)$ , need some def's...

Def: If  $P$  is a tableau, the row word of  $P$  is the perm  $\pi_P = R_l R_{l-1} \dots R_1$  where the  $R_i$  are the rows of  $P$ .

$$\text{Ex: } P = \begin{array}{cccc} 1 & 3 & 5 & 7 \\ 2 & 6 & & \\ 4 & & & \end{array} \Rightarrow \text{row word is } 4 \ 2 \ 6 \ 1 \ 3 \ 5 \ 7$$

Lemma 3: If  $P$  a std tableau, then

$$\pi_P \xrightarrow{RS} (P, -) \quad (\text{Exercise})$$

$$4 \ 2 \ 6 \ 1 \ 3 \ 5 \ 7 \rightarrow \begin{array}{cccc} 1 & 3 & 5 & 7 \\ 2 & 6 & & \\ 4 & & & \end{array}$$

Note: def's & theorems about std tableaux make sense for partial tableaux & partial perms, i.e. bijections  $\pi: K \rightarrow L$  between 2 sets of pos. integers. If  $K = \{k_1 < \dots < k_m\}$ , we can write  $\pi$  as

$$\begin{array}{cccc} k_1 & k_2 & \dots & k_m \\ \downarrow & \downarrow & & \downarrow \\ l_1 & l_2 & \dots & l_m \end{array} \quad \text{where } l_i = \pi(k_i)$$

Can insert the  $l_i$ 's into  $P$  and place  $k_i$ 's in  $Q$ .

Now back to  $(\Leftarrow)$ . By Lemma 3, enough to show that if  $P(\pi) = P$  then  $\pi \stackrel{K}{\cong} \pi_P$ .

Induction on  $n$ . ( $\pi \in S_n$ )

Let  $x$  be last element of  $\pi$ , so  $\pi = \pi' x$   
 and  $\pi'$  a sequence of  $n-1$  letters.

Let  $p' = P(\pi')$ .

By induction,  $\pi' \stackrel{k}{\equiv} \pi_{p'}$

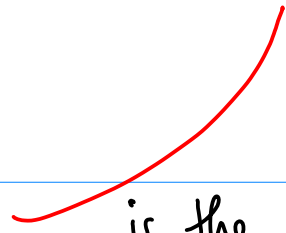

Need to show  $\pi' x \stackrel{k}{\equiv} \pi_p$  so by  $\nearrow$ ,  
 enough to show  $\pi_{p'} x \stackrel{k}{\equiv} \pi_p$ .

Idea: The Knuth relations used to transform  $\pi_{p'} x$  into  $\pi_p$  simulate insertion of  $x$  into  $p'$ .

Let rows of  $p'$  be  $R_1, \dots, R_k$  where  $R_1 = p_1 p_2 \dots p_k$ .  
 If  $x$  enters  $p'$  in column  $j$ , then

$$p_1 < \dots < p_{j-1} < x < p_j < \dots < p_k.$$

$$\begin{aligned} \text{So } \pi_{p'} x &= R_k \dots R_2 p_1 \dots p_{j-1} p_j \dots p_{k-1} p_k x && x < p_{k-1} < p_k \\ &\stackrel{1}{\equiv} R_k \dots R_2 p_1 \dots p_{j-1} p_j \dots p_{k-1} x p_k && x < p_{k-2} < p_{k-1} \\ &\stackrel{2}{\equiv} R_k \dots R_2 p_1 \dots p_{j-1} p_j \dots x p_{k-1} p_k \\ &\vdots \\ &\stackrel{j-1}{\equiv} R_k \dots R_2 p_1 \dots p_{j-1} p_j x p_{j+1} \dots p_k && p_{j-1} < x < p_j \\ &\stackrel{2}{\equiv} R_k \dots R_2 p_1 \dots \overleftarrow{p_j p_{j-1} x} p_{j+1} \dots p_k \\ &\vdots \\ &\stackrel{2}{\equiv} R_k \dots R_2 p_j p_1 \dots p_{j-1} x p_{j+1} \dots p_k && p_{j-2} < p_{j-1} < p_j \end{aligned}$$

Now this  is the first row of  $P = r_x(P')$ .  
 And the element bumped from first row is just  
 after  $R_2$ . So we can continue our  
 sequence of  $\cong^1$  and  $\cong^2$ , until we  
 transform the word into the row word  
 of  $r_x(P') = P$   
 $\circ \circ \quad \pi_{P'} \times \cong^k \pi_P$  

Green generalized Schensted's result on  
 increasing + decreasing sequences.

Def: Let  $\pi$  be a sequence. A subsequence  
 $\sigma$  of  $\pi$  is  $k$ -increasing if, as a set,  
 it can be written as the disjoint union  
 $\sigma = \sigma_1 \cup^* \sigma_2 \cup^* \dots \cup^* \sigma_k$  where the  $\sigma_i$   
 are increasing subsequences of  $\pi$ . If the  $\sigma_i$  are  
 all decreasing, call  $\sigma$   $k$ -decreasing.

Let  $i_k(\pi) =$  length of  $\pi$ 's longest  $k$ -increasing  
 subsequences.  
 $d_k(\pi) =$  " " "  $k$ -decreasing  
 subsequences.

$k=1$  gives usual increasing + decreasing sequences.

Ex:  $\pi = 4 \underline{23651} 7$

Longest 1, 2 and 3-increasing subsequences are:

- 2357
- $4 \ 236517 = 2357 \cup^* 46$
- $4 \ 236517 = 2357 \cup^* 46 \cup^* 1$

So  $i_1(\pi) = 4$ ,  $i_2(\pi) = 6$ ,  $i_3(\pi) = 7$

Recall  $P(\pi) = \begin{array}{c} 1357 \\ 26 \\ 4 \end{array}$

Guess the relation of  $i_j$ 's to shape of  $P(\pi)$ ?

$$\lambda = (4, 2, 1) \quad \text{so} \quad \lambda_1 = 4$$
$$\lambda_1 + \lambda_2 = 6$$
$$\lambda_1 + \lambda_2 + \lambda_3 = 7$$

Theorem (Green): Given  $\pi \in S_n$ , let  
sh  $P(\pi) = (\lambda_1, \dots, \lambda_r)$  w/ conjugate  $(\lambda'_1, \dots, \lambda'_m)$ .

Then for any  $k$ ,

$$i_k(\pi) = \lambda_1 + \dots + \lambda_k,$$
$$d_k(\pi) = \lambda'_1 + \dots + \lambda'_k.$$

Strategy of proof: given a tableau  $P$ ,  
we will prove result for special  
perm  $\pi_P$ . Then show result holds for  
all elems in the equiv class of  $\pi_P$ .  
(all such elems have same shape).



$\pi = \pi_1 \stackrel{k}{\approx} \pi_2 \stackrel{k}{\approx} \dots \stackrel{k}{\approx} \pi_i = \pi_p$  where each  $\stackrel{k}{\approx}$  is given by a Knuth relation.

All  $\pi_i$  have same tableau, so need to show they have same value for  $i_k$ .

Enough to show if  $\pi \stackrel{i}{\approx} \sigma$  for  $i=1$  or  $2$  then  $i_k(\pi) = i_k(\sigma)$ . Let's do case  $\stackrel{1}{\approx}$  and leave  $2$  as exercise.

Suppose  $\pi = x_1 \dots y x z \dots x_n$  and  $\sigma = x_1 \dots y z x \dots x_n$ . To show  $i_k(\pi) \leq i_k(\sigma)$ , show that any  $k$ -incr. subsequence of  $\pi$  has corresp.  $k$ -incr. subsequence in  $\sigma$  of same length.

Let  $\pi' = \pi_1 \cup^* \pi_2 \cup^* \dots \cup^* \pi_k \subset \pi$ . If  $x \neq z$  are not in the same  $\pi_i$ , then the  $\pi_i$  are incr. subsequences of  $\sigma$  so done.

Now suppose  $x \neq z$  both in same  $\pi_i$ , WLOG  $\pi_1$ .

- If  $y \notin \pi'$ , let  $\sigma_1 = \pi_1$  w/  $x$  replaced by  $y$ .  $x < y < z$  so  $\sigma_1$  still increasing &  $\sigma' = \sigma_1 \cup^* \pi_2 \cup^* \dots \cup^* \pi_k$  has same length.

- If  $y \in \pi'$ , say  $y \in \pi_2$  (note: can't have  $y \in \pi_1$ )

Let  $\pi_1' =$  subsequence of  $\pi_1$  up to & including  $x$   
 $\pi_1'' =$  subsequence which is the rest of  $\pi_1$ , } starts w/  $z$   
 $\pi_2' =$  subsequence of  $\pi_2$  up to & including  $y$   
 $\pi_2'' =$  sub " " which is the rest of  $\pi_2$



If we use 29568, no way to complete it.  
Can use 2479... and 1368.

Check w/ class : schedule for December.

Mon	Dec 3	(during section)	in-class presentations
Tues	Dec 4	"	"
Thurs	Dec 6	"	"

Mon	Dec 10	Section :	review for exam
Tues	Dec 11	class :	review for exam
Thurs	Dec 13	no class	

Tues	Dec 18	In-class exam	
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Will everyone be there on Dec 18?

## Sketch of proof of hook-length formula

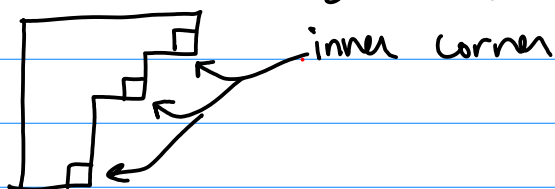
$$\text{Th: } f^\lambda = \frac{n!}{\prod_{v \in \lambda} h_v}$$

Recall  $H_v$  is the hook of  $v$  &  $h_v$  is its length

Probabilistic proof.

Algorithm due to Greene, Nijenhuis, Wilf.

Let  $\lambda$  be Young diagram w/  $n$  boxes.



GNW 1: Pick a box  $v \in \lambda$  w/ prob.  $1/n$ .

GNW 2: While  $v$  is not an inner corner do:

a. Pick a box  $\bar{v} \in H_v - \{v\}$  w/  
probability  $1/(h_v - 1)$

b.  $v := \bar{v}$

GNW 3: Give the label  $n$  to the corner  $v$  you have reached.

GNW 4: Go back to step GNW 1 w/  $\lambda := \lambda - \{v\}$   
and  $n := n - 1$  & repeat outer loop  
until all cells are labeled.

We call the sequence of nodes which we see in one pass thru GWN 1-3 a trial.

Idea: This algorithm produces any given std tableau  $P$  of shape  $\lambda$  w/ probability

$$\text{prob}(P) = \frac{\prod_{v \in \lambda} h_v}{n!} \quad (*)$$

This would prove the theorem!

Step 1: The algorithm terminates + produces a std tableau.

Step 2: Let  $(\alpha, \beta)$  be the box of  $P$  containing  $n$  + let  $\text{prob}(\alpha, \beta)$  be the prob. that the 1st trial ends there. Prove that  $(*)$  follows by induction from

$$+ \text{prob}(\alpha, \beta) = \frac{1}{n} \prod_{i=1}^{\alpha-1} \left(1 + \frac{1}{h_{i,\beta} - 1}\right) \prod_{j=1}^{\beta-1} \left(1 + \frac{1}{h_{\alpha,j} - 1}\right)$$

Certainly knowing this allows us to compute  $\text{prob}(P)$ !

Step 3: Given a trial ending at  $(\alpha, \beta)$ ,  
let the horiz projection of the trial be

$$I = \{i \neq \alpha : v = (i, j) \text{ for some } v \text{ on the trial}\}$$

Similarly define vert projection.

Let  $\text{prob}_{I, J}(\alpha, \beta)$  denote the sum of the prob's of all trials terminating at  $(\alpha, \beta)$  w/ horiz + vert projections  $I \neq J$ .

$$\text{Show: } \text{prob}_{I, J}(\alpha, \beta) = \frac{1}{n} \prod_{i \in I} \frac{1}{h_{i, \beta} - 1} \prod_{j \in J} \frac{1}{h_{\alpha, j} - 1}$$

Use this to show  $\dagger$ .

Induction on  $|I \cup J|$ .

