

155 Lecture 19

The Grassmannian, the complete flag variety, Bruhat order...

The Grassmannian & the complete flag variety are geometric objects that are closely related to the combinatorics (partitions, permutations, symm. functions) that we have studied.

The Grassmannian $Gr_{k,n}(\mathbb{C})$ (or $Gr_{k,n}(\mathbb{R})$) is the set of all k -dim subspaces of \mathbb{C}^n , i.e. $\{V \subset \mathbb{C}^n \mid \dim V = k\}$.

It is a

- cell complex
- manifold
- projective variety

More concrete way of thinking about it:

Can represent $V \in Gr_{k,n}(\mathbb{C})$ as a full rank $k \times n$ matrix; the k rows span a k -dim subspace of \mathbb{C}^n .

But we consider two such matrices A and A' to be equivalent if they span the same subspace, i.e. if

$$\exists P \in GL_k(\mathbb{C}) \text{ s.t. } PA = A'$$

Why does $PA = A'$ say that A and A' span the same subspace?

Write $P \in GL_k(\mathbb{C})$ as $\begin{pmatrix} p_1 \\ \hline p_2 \\ \hline \vdots \\ \hline p_k \end{pmatrix}$ where p_i 's are the row vectors.

Write $A \in \text{Mat}(k, n)$ as $\begin{pmatrix} c_1 & | & c_2 & | & \dots & | & c_n \end{pmatrix}$ where the c_i 's are the column vectors.

Then $PA =$

$$\begin{pmatrix} p_1 \\ \hline p_2 \\ \hline \vdots \\ \hline p_k \end{pmatrix} \begin{pmatrix} c_1 & | & c_2 & | & \dots & | & c_n \end{pmatrix} = \begin{pmatrix} p_1 \cdot c_1 & p_1 \cdot c_2 & \dots & p_1 \cdot c_n \\ p_2 \cdot c_1 & p_2 \cdot c_2 & \dots & p_2 \cdot c_n \\ \vdots & \vdots & \dots & \vdots \\ p_k \cdot c_1 & p_k \cdot c_2 & \dots & p_k \cdot c_n \end{pmatrix}$$

↙ dot product

On the RHS, note that each row, say $(p_i \cdot c_1 \quad p_i \cdot c_2 \quad \dots \quad p_i \cdot c_n)$ is just a

linear combination of the rows of $\begin{pmatrix} c_1 & | & c_2 & | & \dots & | & c_n \end{pmatrix}$

∴ the span of the rows of the RHS is contained in the span of the rows of $\begin{pmatrix} c_1 & | & c_2 & | & \dots & | & c_n \end{pmatrix}$

Since P is invertible, the two row spans are the same.

◦◦ we can identify $Gr_{kn}(\mathbb{C})$ with
 $GL_k(\mathbb{C}) \setminus Mat(k, n)$ \swarrow full rank $k \times n$ matrices

Example: $Gr_n(\mathbb{C})$.

The set of 1-planes (lines) in \mathbb{C}^n , i.e. projective space.

The set of full rank $1 \times n$ matrices s.t.

$(a_1, \dots, a_n) \sim (b_1, \dots, b_n)$ if $\exists \lambda \in GL_1(\mathbb{C}) = \mathbb{C}^* \text{ s.t.}$
 $\lambda(a_1, \dots, a_n) = (b_1, \dots, b_n)$.

In general the topology of $Gr_{kn}(\mathbb{C})$ is complicated. But we can partition $Gr_{kn}(\mathbb{C})$ into smaller pieces (actually, cells) which are easy to understand. This will give $Gr_{kn}(\mathbb{C})$ the structure of a cell complex.

Take any matrix A representing an element of $Gr_{kn}(\mathbb{C})$.

Note that performing row operations on A leaves the corresp. element of $Gr_{kn}(\mathbb{C})$ unchanged.

Algorithm: ◦ Look at each row & locate the last nonzero entry. Then permute rows, ordering them by the position of that nonzero entry.

Each * represents an element of \mathbb{F}

Get something like:

$$\begin{pmatrix} * & * & * & * & * & * & * \\ * & * & * & * & * & 0 & 0 \\ * & * & * & * & * & 0 & 0 \\ * & * & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

- Within each row, divide every entry by the last nonzero entry

Get something like:

$$\begin{pmatrix} * & * & * & * & * & * & 1 \\ * & * & * & * & 1 & 0 & 0 \\ * & * & * & * & 1 & 0 & 0 \\ * & 1 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

- If 2 rows r_i & r_j both have the initial 1 in the same position, replace r_j w/ $r_j - r_i$.
(& divide again by last nonzero entry)

Get something like:

$$\begin{pmatrix} * & * & * & * & * & * & 1 \\ * & * & * & * & 1 & 0 & 0 \\ * & * & * & 1 & 0 & 0 & 0 \\ * & 1 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

- For i from 1 to k , take row r_i & subtract from it appropriate multiples of rows r_j ($i < j \leq k$) in order to put $k-j$ 0's into the new row r_i .

Get something like:

$$\begin{pmatrix} * & 0 & * & 0 & 0 & * & 1 \\ * & 0 & * & 0 & 1 & 0 & 0 \\ * & 0 & * & 1 & 0 & 0 & 0 \\ * & 1 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

Deleting the identity submatrix from above we get an array of complex numbers:

$$\begin{array}{c} * * * \\ * * \\ * * \\ * \end{array}$$

Their shape represents a partition (here, $(3, 2, 2, 1)$).

Observation: This algorithm associates to any $A \in Gr_{k,n}(\mathbb{C})$ a unique matrix of the form

where $*$'s represent elements of \mathbb{C} .

$$\begin{pmatrix} * & 0 & * & 0 & 0 & * & 1 \\ * & 0 & * & 0 & 1 & 0 & 0 \\ * & 0 & * & 1 & 0 & 0 & 0 \\ * & 1 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

The array of $*$'s forms a partition λ whose Young diagram $\subseteq k \times (n-k)$ rectangle.

All such 0-1- $*$ matrices arise in this way for all Young diagrams $\lambda \subseteq k \times (n-k)$.

Def: If we fix λ — equivalently, fix the form of the matrix as in

$$\begin{pmatrix} * & 0 & * & 0 & 0 & * & 1 \\ * & 0 & * & 0 & 1 & 0 & 0 \\ * & 0 & * & 1 & 0 & 0 & 0 \\ * & 1 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

The set of all elements of $Gr_{k,n}(\mathbb{C})$ of this form is called the Schubert cell Ω_λ

Topologically this is just \mathbb{C}^d where $d = |\lambda|$ (# of boxes of λ).

$$\text{So } Gr_{k,n}(\mathbb{C}) = \bigsqcup_{\lambda \leq k \times (n-k)}^* \Omega_\lambda \quad (\text{disjoint union})$$

A natural question is: when is one cell $\Omega_\lambda \subset \Omega_\mu$?

Answer: $\Omega_\lambda \subset \Omega_\mu$ iff $\lambda \leq \mu$. (Young diagrams) (containment of)

Fact (beyond scope of course): The cohomology ring $H^*(Gr_{k,n}(\mathbb{C}))$ is a polynomial ring which is generated by the Schur functions S_λ for all $\lambda \leq k \times (n-k)$

Now: Flag varieties.

Fix n and consider a subset $\{i_1 < i_2 < \dots < i_r\} \subset \{1, \dots, n\}$.

We define the partial flag variety $Fl(i_1, \dots, i_r)$ to be the set of flags in \mathbb{C}^n

$$\{V_1 \subset V_2 \subset \dots \subset V_r \subset \mathbb{C}^n \mid \dim V_j = i_j\}.$$

↖ each V_i a vector space.

Each chain of subspaces $V_1 \subset \dots \subset V_r$ is called a flag.

Note that if $\{i_1 < \dots < i_r\} = \{k\}$, the corresponding flag variety is just the Grassmannian $Gr_{k,n}(\mathbb{C})$.

When $\{i_1 < \dots < i_r\} = \{1, \dots, n\}$, this is called the complete flag variety denoted Fl_n .

Just as we did for the Grassmannian, we can represent a flag variety by a matrix. Eg if A is a full rank $n \times n$ matrix, w/ rows

$$\begin{pmatrix} \hline a_1 \\ \hline a_2 \\ \hline \vdots \\ \hline a_n \end{pmatrix},$$

then this corresponds to the complete flag $(\text{span}(a_1), \text{span}(a_1, a_2), \dots, \text{span}(a_1, a_2, \dots, a_n))$.

As w/ the Grassmannian, we can perform certain operations on the matrix which preserve the corresp element of Fl_n .

[Ask what operations are legit!]

We can:

- multiply a row by a scalar

- If $i < j$, we can replace r_j with $r_j - \lambda r_i$.

But we can't do this if $i > j$ (why?)

And we can't permute rows (why?)

So there is again an algorithm by which we

can take a matrix $\begin{pmatrix} a_1 \\ \vdots \\ a_n \end{pmatrix}$ representing an element of $F[x]$ & put it in a canonical form.

Algorithm:

- Divide each row by its last nonzero entry, getting something like

Again, α 's represent entries in \mathbb{C} \rightarrow

$$\begin{pmatrix} \alpha & \alpha & \alpha & \alpha & 1 & 0 \\ \alpha & 1 & 0 & 0 & 0 & 0 \\ \alpha & \alpha & \alpha & \alpha & \alpha & 1 \\ \alpha & \alpha & \alpha & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 \\ \alpha & \alpha & 1 & 0 & 0 & 0 \end{pmatrix}$$

- Take row r_j & subtract off appropriate multiples of r_i (for all $i < j$) to put 0's in as many entries as possible of r_j .

Get something like

$$\begin{pmatrix} \alpha & \alpha & \alpha & \alpha & 1 & 0 \\ \alpha & 1 & 0 & 0 & 0 & 0 \\ \alpha & 0 & \alpha & \alpha & 0 & 1 \\ \alpha & 0 & \alpha & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \end{pmatrix}$$

Notice that the positions of the 1's defines

a permutation π . Here, it is $(5, 2, 6, 4, 1, 3)$

Note that the $*$'s which remain exactly correspond to entries of the matrix w/ which have a 1 to the right & a 1 below. Letting $i \neq j$ be the rows of those 1's ($i < j$), these correspond to when $\pi(i) > \pi(j)$.

Def: An inversion of a permutation π is a pair $i < j$ s.t. $\pi(i) > \pi(j)$.

Observation: Each element of F_n can be uniquely represented in the form

$$\begin{pmatrix} * & * & * & * & 1 & 0 \\ * & 1 & 0 & 0 & 0 & 0 \\ * & 0 & * & * & 0 & 1 \\ * & 0 & * & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \end{pmatrix}$$

for some $\pi \in S_n$ & where each $*$ is replaced w/ an element of \mathbb{C} .

And all matrices of this form represent elements of F_n .

If we fix $\pi \in S_n$, the set of all elements of F_n represented by

where position of 1's gives π \rightarrow

$$\begin{pmatrix} * & * & * & * & 1 & 0 \\ * & 1 & 0 & 0 & 0 & 0 \\ * & 0 & * & * & 0 & 1 \\ * & 0 & * & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \end{pmatrix}$$

is called the Schubert cell Ω_π .
Topologically, $\Omega_\pi \cong \mathbb{C}^{\#\text{inv}(\pi)}$

Recall S_n is generated by $\langle s_i \mid 1 \leq i \leq n-1 \rangle$
where $s_i = (i \ i+1)$

Def: The length $l(\pi)$ of $\pi \in S_n$ is
the minimal number r for which we
can write $\pi = s_{i_1} s_{i_2} \dots s_{i_r}$.
Any such minimal factorization of π is
called a reduced expression of π .

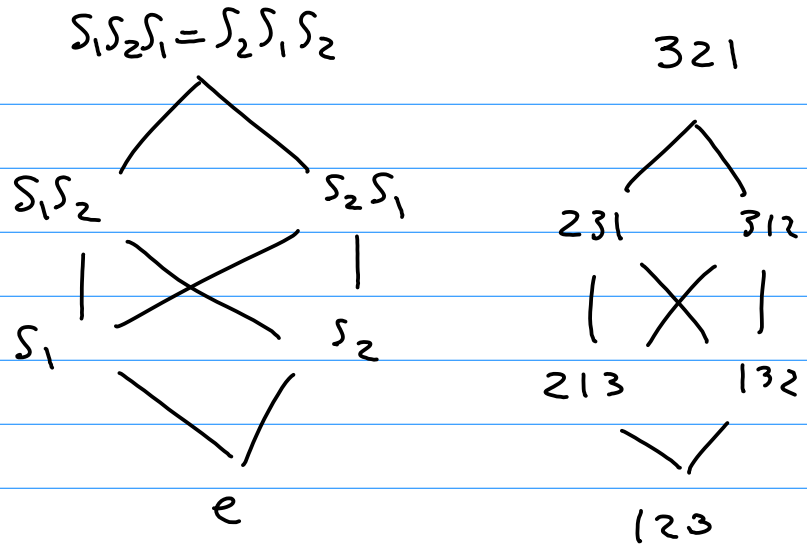
In your homework, you proved (hopefully)
that $l(\pi) = \#\text{inversions of } \pi$.

There is more structure on S_n that is
related to $l(\pi)$: S_n has the
structure of a partially ordered set
(poset), under the Burhat order.

Def: Let $s_{i_1} \dots s_{i_r}$ be a reduced decomp.
of a permutation w . Then we say
 $v \leq w$ iff \exists a subsequence
 (j_1, \dots, j_m) of (i_1, \dots, i_r) s.t.
 $s_{j_1} \dots s_{j_m}$ is a reduced decomp. of v

This is the Burhat order. (also known as
strong order or
strong Burhat order)

For S_3 :



This depiction of the poset is called the Hasse diagram.

We put edges between elements of the poset that are cover relations, i.e. $u < v$ where $\nexists w$ s.t. $u < w < v$.

In this example, there is a rank function on the poset. i.e. since all maximal chains have the same length, we can define $\text{rank}(w)$ to be the # of elements in a maximal chain $e < w^1 < w^2 < \dots < w^r = w$.

And for the Bruhat order, $\text{rank}(w) = \ell(w)$.

Question: When do we have $\Omega_v \subset \Omega_w$? (Schubert cells)

Answer: Exactly when $v \leq w$ in the Bruhat order.