

Problem 1

PS2, 1

The \pm -moves $\uparrow \rightarrow \downarrow$ and $\uparrow \rightarrow \downarrow$ change the writhe by ± 1 w/o affecting the link. We apply them to any given diagram to get $w(L) = 0$.

Problem 2 [Tim]

If K_1 & K_2 are the two link components, color the projection of K_1 using the chessboard coloring. Then at each crossing between K_1 & K_2 , the strand in K_2 goes from a black square to a white one, or viceversa. Since we must start & end in the same color when we travel once around K_2 , we must have an even # of crossings between K_1 & K_2 . Each of these crossings adds 1 mod 2 to the signed sum of crossings, which then turns out to be even. Hence the linking number is even, and we are done.



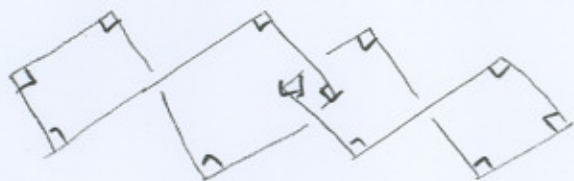
Problem #3



Assume we start travelling from \vec{v} shown in the diagram to the right. \vec{v} is contained in the plane determined by e_1 & e_2 so, as it travels from e_1 to e_2 , its unitary representative centered at the origin traces a great circle arc on S^2 . As \vec{v} travels around the knot and returns to its startpoint, its representative traces a closed curve on S^2 . The winding # of the knot is then $\frac{\Delta\theta}{2\pi}$ times the angle change of the representative vector. But since we traced a closed curve of great circle arcs, this # must be an integer.

Problem 4

For links, the answer is yes!



$$\begin{aligned} \text{writhe} &= 0 \\ \text{winding} &\neq 0 \end{aligned}$$

for knots, the answer is no!

In general, the difference (winding # - writhe) is $1 \pmod{2}$

for all knots:

Note that we can unknot a knot by changing crossings in any of its diagrams, and each switch changes writhe by 2. Furthermore, 1-moves change both writhe + w.n. by 1 and R_2, R_3 leave both unaffected. We can use these moves to transform any knot diagram into a planar S^1 , which has w.n.=1 and writhe 0, this proves the claim.

Problem 5

$$\cdot \quad \begin{array}{c} \diagup \\ | \\ \diagdown \end{array} \rightarrow \begin{array}{c} \circ \\ | \\ \circ \end{array} \rightarrow \begin{array}{c} \circ \\ \bigcirc \\ | \\ \circ \end{array} \rightarrow \begin{array}{c} \circ \\ \bigcirc \\ | \\ \circ \end{array} \rightarrow \begin{array}{c} \circ \\ | \\ \circ \end{array} \rightarrow | \end{array}$$

$$\cdot \quad \begin{array}{c} \diagup \\ \diagdown \end{array} \rightarrow \begin{array}{c} \diagup \\ \bigcirc \\ \diagdown \end{array} \rightarrow \begin{array}{c} \diagup \\ \bigcirc \\ \bigcirc \\ \diagdown \end{array} \rightarrow \begin{array}{c} \diagup \\ \diagdown \end{array}$$



Problem 6

PS2, 4

Hopf



→ black surface:



orientable!

→ white surface:



orientable!

[black here just to illustrate... not the same black as before]

Trefoil



→ black



unorientable!

(odd # of twists)

→ white: orientable

remove ∞ from bottom disk



Figure 8



→ black: unorientable



$0 \pmod 2$

→ white: orientable

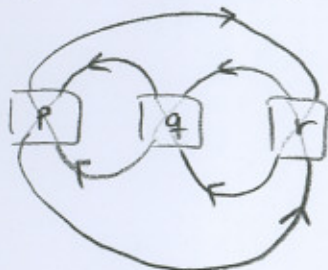


Seifert surface!



Problem 7

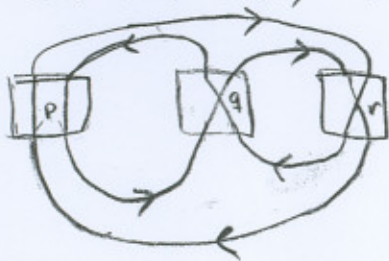
If p, q, r are all odd, the knot may be oriented as



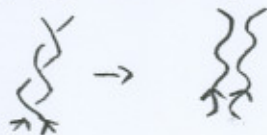
so the Seifert algorithm produces $(p+1) + (q+1) + (r+1) + 2$ disks. Since there are $p+q+r$ crossings and $\mu(D) = 1$,

$$\text{the genus is } \frac{1}{2} (1 - \underbrace{(p+1) + (q+1) + (r+1)}_1 + 2) + (p+q+r+1-1)$$

From pset 1, the case $(p, q, r) = (0, 1, 1) \pmod{2}$ suffices to show what happens when only one of the (p, q, r) triplet is even. In this case, the orientation is given by the diagram



Then the $q+r$ produces no Seifert circles:



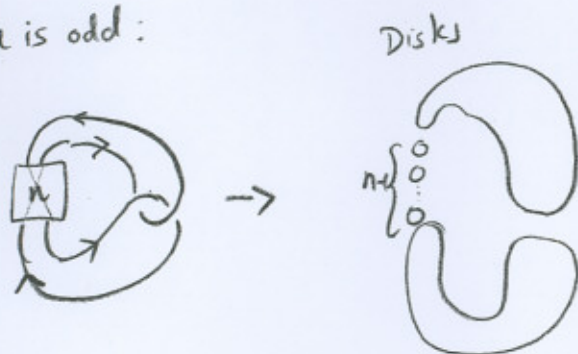
and hence the genus is

$$\frac{1}{2} (1 - p - 1 + (p+q+r)) = \frac{q+r}{2}$$

□

Problems

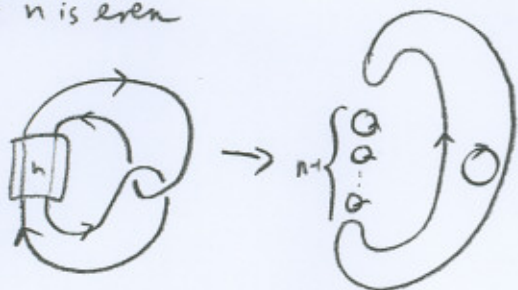
• If n is odd:



we have $n+2$ crossings and $(n-1)+2 = n+1$ disks, so
the genus is

$$\frac{1}{2} (1 - (n+1) + (n+2)) = 1$$

• If n is even



We have $n+2$ crossings and $n+1$ disks, so
by the same computations above, the genus is 1.



[Diagrams by Eleanor]

